

- \*S63  $(P \rightarrow Q) \rightarrow R, S \rightarrow (\sim Q \rightarrow T) \vdash R \vee \sim T \rightarrow (S \rightarrow R)$   
 \*S64  $(P \rightarrow Q) \& (R \rightarrow P), (P \vee R) \& \sim(Q \& R) \vdash (P \& Q) \& \sim R$   
 S65  $P \& Q \rightarrow (R \vee S) \& \sim(R \& S), R \& Q \rightarrow S,$   
 $S \rightarrow ((R \& Q) \vee (\sim R \& \sim Q)) \vee \sim P \vdash P \rightarrow \sim Q$   
 S66  $\sim(P \& \sim Q) \vee \sim(\sim R \& \sim S),$   
 $\sim S \& \sim Q, T \rightarrow (\sim S \rightarrow \sim R \& P) \vdash \sim T$   
 \*S67  $P \& Q \rightarrow R \vee S \vdash (P \rightarrow R) \vee (Q \rightarrow S)$   
 \*S68  $P \& Q \rightarrow (R \vee S) \& \sim(R \& S), R \& Q \rightarrow S,$   
 $S \rightarrow ((R \& Q) \vee (\sim R \& \sim Q)) \vee \sim P \vdash P \rightarrow \sim Q$   
 S69  $(P \rightarrow Q) \rightarrow (Q \rightarrow P) \vdash (P \rightarrow Q) \rightarrow (\sim P \rightarrow \sim Q)$

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 2.3 Theorems
**Theorem**

*Definition.* A **THEOREM** is a sentence that can be proved from the empty set of premises.

*Comment.* We can assert that a given sentence is a theorem by presenting it as the conclusion of a sequent with nothing to the left of the turnstile.

Example.

Prove  $\vdash P \& Q \rightarrow Q \& P$ .

1	(1)	$P \& Q$	A
1	(2)	$Q$	1 &E
1	(3)	$P$	1 &E
1	(4)	$Q \& P$	2,3 &I
	(5)	$P \& Q \rightarrow Q \& P$	4 $\rightarrow$ I(1)

*Comment.* Note that in step 5 we discharge assumption 1. Hence, the final conclusion rests on no assumptions (i.e., the assumption-set is the empty set).